'An order is to be obsenved in all things'

## A STUDY ON SEQUENCING MODELS

## BY

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Definition: Selection of an appropriate order for a series of job to be done on a number of service facilities so as to optimize the total effectiveness (may be time, cost etc which is a function of the order)

Number of machines
Processing order Processing time Idle time on a machine
Total elapsed time No Passing Rule


## =Machine operations in Fothe



- Machine order fixed - Turning Threading Knurling Job order 123456132456 . . .

234561243561 ...
345612354612 ...
456123465123 . . .
561234516234 . . .
612345621345 . . . Total 6! Orders

- From these find the order which minimizes time/ cost
This is a 6 job - 3 machine problem



## Sequencing in computer systems



- $n$ jobs on $m$ machines
(n!) ${ }^{m}$ possible sequences
Find the sequence minimizing the total time When $\mathrm{n}=4, \mathrm{~m}=5$ there are
$(4!)^{5}=7962624$ possible sequences
Enumeration impossible for even smaller $m$ and n


## $\equiv$ Processing n jobs through ${ }^{2} \equiv$ machines $=$

General Form

$$
\begin{array}{l|cccc}
\text { Jobs } & 1 & 2 & \ldots & n \\
\hline \operatorname{Mac} A & a_{1} & a_{2} & \ldots & a_{n} \\
\operatorname{Mac} B & b_{1} & b_{2} & \ldots & b_{n}
\end{array}
$$



## Gnatt Chart

Consider the two job Two machine problem


| Job | $\mathrm{J}_{1}$ | $\mathrm{~J}_{2}$ |
| :--- | :---: | :---: |
| Machine A | 3 | 5 |
| Machine B | 5 | 4 |


|  | $\mathrm{J}_{1}(3)$ |  |
| :---: | :---: | :---: |
| A | $\mathrm{J}_{2}(4)$ | $\mathrm{J}_{1}(5)$ |

Total Time $=14$

Illustration:A book binder has a printing press, a binding machine \& the manuscripts of different books. The processing times are given. Determine the optimum sequence.

| BOOK | 1 | 2 | 3 | 4 | 5 | 6 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Printing <br> Time (hrs) | 30 | 120 | 50 | 20 | 90 | 100 |
| Binding <br> Time (hrs) | 80 | 100 | 90 | 60 | 30 | 10 |

## Eptifinu sequence Algofithm

Find the smallest processing time. If it is for first machine, place the job in first available position of the sequence. If it is for second machine place the job in the last available position of the sequence. If there is a tie


## OOptimum sequence Algoftitn =

(i) among the two machines ( $a_{k}=b_{r}$ ) place job corresponding to first machine (job k) in first available position of the sequence and the job corresponding to second machine (job r) in the last available position of the sequence.
(ii) among same machine $\left(a_{k}=a_{r} / b_{k}=b_{r}\right)$, break the tie arbitrarily


Fptimum sequence Algofithem
Cross off the jobs already assigned and repeat the above procedure


Calculate the idle times and total elapsed time

| Books | Printing <br> In <br> Out | Binding Out <br> In |  | Idle times |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Printing | Binding |  |  |  |  |
| 4 | $0-20$ | $20-80$ | 0 | 20 |  |
| 1 | $20-50$ | $80-160$ | 0 | 0 |  |
| 3 | $50-100$ | $160-250$ | 0 | 0 |  |
| 2 | $100-220$ | $250-350$ | 0 | 0 |  |
| 5 | $220-310$ | $350-380$ | 0 | 0 |  |
| 6 | $310-410$ | $410-420$ | 10 | 30 |  |

General form

| Jobs | 1 | 2 | $\ldots$ | $n$ |
| :--- | :--- | :--- | :--- | :--- |
| Mac A | $a_{1}$ | $a_{2}$ | $\ldots$ | $a_{n}$ |
| Mac B | $b_{1}$ | $b_{2}$ | $\ldots$ | $b_{n}$ |
| Mac C | $c_{1}$ | $c_{2}$ | $\ldots$ | $C_{n}$ |

- Solution Procedure

If minimum among $a_{i} /$ among $c_{c}$ is greater than or equal to the maximum among $c$; then we can reduce this to an $n$ job -2 machine problem as below:

Introduce two fictitious machines G and H whose processing times are defined by

$$
g_{i}=a_{i}+b_{i} \& h_{i}=b_{i}+c_{i}
$$

Then proceed as in the above case

## EProsessing $\bar{n}$ jobs throught

 machines- General form

| Jobs | 1 | 2 | $\ldots$ | $n$ |
| :--- | :--- | :--- | :--- | :--- |
| Mac $A$ | $a_{1}$ | $a_{2}$ | $\ldots$ | $a_{n}$ |
| Mac B | $b_{1}$ | $b_{2}$ | $\ldots$ | $b_{n}$ |
| $\ldots$ | $\ldots$ | $\ldots$ |  |  |

Mac $F \quad f_{1} \quad f_{2} \ldots f_{n}$

- Solution Procedure

If minimum among $a_{i}$ / among $f_{i}$ is greater than or equal to the maximum among $b_{i}, c_{i}, \ldots, e_{i}$ then we can reduce this to an $n$ job -2 machine problem as below


## Introduce two fictitious machines G

 and H whose processing times are defined by$\mathrm{g}_{\mathrm{i}}=\mathrm{a}_{\mathrm{i}}+\mathrm{b}_{\mathrm{i}}+\ldots+\mathrm{e}_{\mathrm{i}} \&$
$h_{i}=b_{i}+c_{i}+\ldots+f i$
Then proceed as in the above case
 4 job - 6 machine problem.

$$
\begin{array}{l|rllllll}
\text { Machines } & M_{1} & M_{2} & M_{3} & M_{4} & M_{5} & M_{6} \\
\hline \text { Job A } & 18 & 8 & 7 & 2 & 10 & 10 & \\
\text { Job B } & 17 & 6 & 9 & 6 & 8 & 19 & \\
\text { Job C } & 11 & 5 & 8 & 5 & 7 & 15 & \\
\text { Job D } & 20 & 4 & 3 & 4 & 8 & 12 & \\
\hline
\end{array}
$$

Min of $M_{1}=11, \operatorname{Max} M_{2}, M_{3 \prime}, M_{4^{\prime}} M_{5}=10$, Min $M_{6}=12$

## Condition for fictitious Machines is satisfied.

Jobs A B C D

Fictitious Machine G 45463639
Machine H $\quad 524840 \quad 31$
Sequence


| Jobs | $M_{1}$ <br> In Out | $M_{2}$ <br> In Out | $M_{3}$ <br> In Out | $M_{4}$ <br> In Out | $M_{5}$ <br> In Out | $M_{6}$ <br> In Out |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: |
| C | $0-11$ | $11-16$ | $16-24$ | $24-29$ | $29-36$ | $36-51$ |
| A | $11-29$ | $29-37$ | $37-44$ | $44-46$ | $46-56$ | $56-81$ |
| B | $29-46$ | $46-52$ | $52-61$ | $61-67$ | $67-75$ | $81-100$ |
| D | $46-66$ | $66-70$ | $70-73$ | $73-77$ | $77-85$ | $100-112$ |

## Total elapsed time $=112 \mathrm{hrs}$

Idle times M1 to M6 are 46, 89, 85, 95,79,41 respectively


